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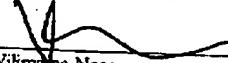
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PHILIPS**Philips Intellectual Property & Standards****FACSIMILE TRANSMITTAL TO
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TO:	FACSIMILE NO.:	TELEPHONE NO.:
Mail Stop Amendment Commissioner for Patents P.O. Box 1450 Arlington, VA 22313-1450	(571) 273-8300	() -
ATTENTION:	<i>Examiner: PAN DANIEL H.</i>	
	<i>Art Unit: 2183</i>	
FROM:	TELEPHONE NO.:	
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RE:	<i>Serial No.: 09/801,080</i>	
	<i>Attorney Docket No.: NL000133</i>	

TRANSMISSION INCLUDES:**13 Pages (including cover sheet)**Transmittal of Brief in Support of Appeal - 1 pageAppeal Brief - 11 pages

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on 8/8/06by 
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IN THE UNITED STATES PATENT AND TRADEMARK OFFICE
BEFORE THE BOARD OF PATENT APPEALS AND INTERFERENCES

First-Named Inventor: BUSA, Natalino Giorgio
Application No.: 09/801,080 Conf.: 5082
Date Filed: 03/07/2001
Customer No.: 24738

Atty Docket No.: NL000133
Art Unit: 2183
Examiner: PAN, DANIEL H

Title: Data Processing Device, Method of Operating a Data Processing Device and Method for Compiling a Program

Mail Stop Appeal Brief-Patents
Commissioner for Patents
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Alexandria, VA 22313-1450

**TRANSMITTAL OF
BRIEF IN SUPPORT OF AN APPEAL**

Sir:

Enclosed is an Appeal Brief in the above-identified patent application.

Please charge the any and all required fees to Deposit Account No. 14-1270.

Date: AUG - 8 - 2006

Respectfully submitted,
PHILIPS ELECTRONICS NORTH AMERICAN CORP.

By

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IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

Before the Board of Patent Appeals and Interferences

In re the Application

Inventor : Busa
Application No. : 09/801,080
Filed : 03/07/2001
For : Data processing device, method of operating a data processing device and method for compiling a program

APPEAL BRIEF

On Appeal from Group Art Unit 2183

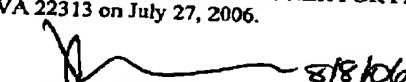
Date: July 27, 2006

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Vilimaina Naga
(Name)


(Signature and Date) 8/8/06

APPEAL
Serial No.: 10/801,080

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RELATED PROCEEDINGS

EVIDENCE

TABLE OF CASES

NONE

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Serial No.: 10/801,080

I. REAL PARTY IN INTEREST

The real party in interest is the assignee of the present application, Koninklijke Philips Electronics N.V., and not the party named in the above caption.

II. RELATED APPEALS AND INTERFERENCES

With regard to identifying by number and filing date all other appeals or interferences known to Appellant which will directly effect or be directly affected by or have a bearing on the Board's decision in this appeal, Appellant is not aware of any such appeals or interferences.

III. STATUS OF CLAIMS

Claims 1-5 are pending, stand finally rejected, and forms the subject matter of the present appeal.

IV. STATUS OF AMENDMENTS

All amendments have been entered. No amendment after final rejection has been submitted.

V. SUMMARY of the CLAIMED SUBJECT MATTER

The present invention relates, in one aspect, to a computational method in which one functional unit, during execution of an instruction, outputs data to or receives input data from another functional unit *during execution of the instruction*, where the functional units share a common memory, e.g., a micro code memory. Claims 1, 4 and 6 have been amended to make clear that such input/output occurs during execution of the

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instruction. Figures 6a and 6b illustrate a particular computation (that of Figure 5, which in turn references the 2Dtransform operation of Figure 2) performed conventionally (Figure 6a) and using the computational method of the present invention (Figure 6b). In the illustrated example, the computational method of the invention results in a 20% reduction in computation time.

The foregoing example is described more particularly in the present specification at page 10, line 27, to page 6, line 19. In Figure 6a, the 2Dtransform operation (Figure 2) is performed *atomically*; i.e., inputs are presented simultaneously to a complex functional unit, which performs the operation and presents outputs simultaneously. In Figure 6b, by contrast, inputs are not presented simultaneously, nor are outputs presented simultaneously. The input $i_1 (= p + q)$ is presented to the complex functional unit as soon as it is available, in cycle 1. During cycles 1 and 2, the second input $i_2 (= p - q - 2)$ is prepared and is then presented to the complex functional unit in cycle 3. Also in cycle 3, the first output $o_1 (= 2 \cdot i_2 + 3)$ is presented. During cycles 3, 4 and 5, the output $o_2 (= 5 \cdot i_1 + 2 \cdot i_2 + 1)$ is computed and presented. During cycles 4 and 5, o_1 (previously made available in cycle 3) is squared and the value 100 is subtracted to form a partial result as part of the condition checking in the last statement of Figure 5. In cycles 6 and 7, o_1 is squared and the resulting quantity added to the partial result. Finally, in cycle 8, the inequality is evaluated. In this example, it may be seen that output data (e.g., o_1) of the complex functional unit ("first functional unit") is processed by the other functional unit ("second functional unit") during execution of the instruction (e.g., 2Dtransform) by the complex functional unit ("first functional unit") as recited in claim 1. Also, input data (e.g., i_2) to the first functional unit is generated by the second functional unit during execution of the instruction as recited in claim 1.

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VI. GROUNDs of REJECTION to be REVIEWED ON APPEAL

The issues in the present matter are whether:

1. claims 1-5 are anticipated by O'Connor.
2. claims 1-5 are anticipated by McNeill.

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Serial No.: 10/801,080**VII. ARGUMENT****I. Rejection of Claims 1-5 as anticipated by O'Connor**

The O'Connor reference teaches something quite different. O'Connor relates to "scoreboarding," a method of resource arbitration in which a centralized set of tables ("registers") in the CPU keep track of what is being used and when. Each row indicates what is being used at each clock tick. Scoreboarding is used to manage dispatch, stalling, and completion of instructions, to watches for Write-After-Write, Read-After-Write and Write-After-Read (WAW, RAW, WAR) hazards, and to manage the execution sequence (e.g. possible instruction reordering) to avoid these hazards. Hazards are detected by observing register references and operation types.

In particular, O'Connor describes a hardware-efficient implementation of a bypass circuit that enables a result from one execution unit to be provided to a waiting execution unit at the same time that result is sent to a register.

In O'Connor, one execution unit waits for results from another execution unit in order for the first execution unit to begin execution. The execution units do not operate concurrently to achieve execution of an instruction.

II. Rejection of Claims 1-5 as anticipated by McNeill

The rejection based on McNeill ignores the basic distinction between an *instruction* and a computing *task*. An instruction is a well-defined operation that can